# Efficient Language Modeling Algorithms with Applications to Statistical Machine Translation 

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## CPU and RAM Costs Matter

"had to favor speed over performance" [Moreau et al, 2013]
"could not test whether this result holds in a large scale evaluation" [Durrani et al, 2013]

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"could not test whether this result holds in a large scale evaluation" [Durrani et al, 2013]

## 0.5-1.9\% BLEU gain from English Gigaword


[Koehn et al, 2012]

## Application: Syntactic MT



## Application: Syntactic MT



## Speech Recognition



Some of the thesis is already used for speech [Kim et al, 2012; Si et al, 2013].

## Optical Character Recognition



Numen (2013) is using most of this thesis for OCR.

## Language Models Are Expensive

Store a sparse set of 121 billion $n$-grams

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$\Longrightarrow$ RAM
$\Longrightarrow \mathrm{CPU}$

## Language Models Are Expensive

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## $\Longrightarrow$ RAM

Millions of probability queries per sentence

Probability does not multiply when strings are concatenated:

$$
\begin{gathered}
p(\text { saw the man }) \\
\neq \\
p(\text { saw }) p(\text { the man })
\end{gathered}
$$

$\Longrightarrow$ Search is hard $\Longrightarrow$ CPU

## Thesis Problem

Much of the CPU and RAM cost is due to the language model. Researchers routinely compromise quality due to these costs.

## Costs Due To Language Models

Estimation from text
Probability queries
Search when the objective includes log probability

## Results Preview

## Speed RAM

Estimation from Text 7.1x 0.07x Raw Queries 2.4x 0.57x Decoding 3.2-10.0x 0.85x

Decoding performance includes $\approx 1.15 x$ speedup from raw queries.
Baseline: SRILM and cube pruning (more later).

## Outline

## Speed RAM Published

| 1 Estimation from Tex | 7.1x | 0.07x | ACL 2013 |
| :---: | :---: | :---: | :---: |
| 2 Raw Queries | $2.4 \times$ | 0.57x | WMT 2011 |
|  |  |  | (IWSLT 2011, |
| 3 Decoding | 3.2-10.0x | 0.85x | EMNLP 2012, |
|  |  |  | NAACL 2013 |

# Estimating LMs is Costly 

MIT RAM
SRI RAM, time
IRST RAM, time, approximation
Berkeley RAM, time, approximation

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Microsoft Delay some computation to query time Google 100-1500 machines, optional stupid backoff

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Microsoft Delay some computation to query time Google 100-1500 machines, optional stupid backoff
"When, oh when, will there be an alternative?"

## Implz Features

- Disk-based streaming and sorting
- User-specified RAM
- Fast
- Interpolated modified Kneser-Ney
$7.7 \%$ of SRI's RAM, $14 \%$ of SRI's wall time




## Adjusting

Adjusted counts are:
Trigrams Same as counts.
Others Number of unique words to the left.

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Adjusted counts are:
Trigrams Same as counts.
Others Number of unique words to the left.

| $\begin{array}{cccr}\text { Suffix Sorted Input } \\ 3 & 2 & 1 & \text { Count }\end{array}$ |  | Output | Output |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | 1-gram Adjusted | 2-gram | Adjusted |
| are one of | 1 | of 2 | one of | 2 |
| is one of | 5 |  | two of | 1 |
| are two of | 3 |  |  |  |

## Streaming Framework

Memory is divided into blocks. Blocks are recycled.


Prepare for next step.

## Adjusted Counts Detail



Each vertex is a thread $\Longrightarrow$ Simultaneous disk and CPU.

## Experiment: Toolkit Comparison

Task Build an unpruned 5-gram LM
Data Subset of English ClueWeb09 (webpages)
Machine 64 GB RAM
Output Format Binary (or ARPA when faster)

IRST disk: 3-way split. Peak RAM of any one process (as if run serially).
Berkeley: Binary search for minimum JVM memory.





## Scaling

## Tokens Smoothing Machines Days

## This Work <br> 126 billion <br> Kneser-Ney <br> 1 2.8

## Counts

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ |
| ---: | :---: | :---: | :---: | :---: | :---: |
| This Work 126B | 393 m | $3,775 \mathrm{~m}$ | $17,629 \mathrm{~m}$ | $39,919 \mathrm{~m}$ | $59,794 \mathrm{~m}$ |
| Pruned Google 1T | 14 m | 315 m | 977 m | $1,313 \mathrm{~m}$ | $1,176 \mathrm{~m}$ |

(This work used a machine with 140 GB RAM and a RAID5 array.)

## Scaling

## Tokens Smoothing Machines Days Year

## This Work

 Google Google 230 billion Google 1800 billion126 billion
31 billion
Kneser-Ney
Kneser-Ney
Kneser-Ney
Stupid

| $\mathbf{1}$ | 2.8 | 2013 |
| ---: | ---: | ---: |
| 400 | 2 | 2007 |
| $?$ | $?$ | 2013 |
| 1500 | 1 | 2007 |

## Counts

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ |
| ---: | :---: | :---: | :---: | :---: | :---: |
| This Work 126B | 393 m | $3,775 \mathrm{~m}$ | $17,629 \mathrm{~m}$ | $39,919 \mathrm{~m}$ | $59,794 \mathrm{~m}$ |
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## Workshop on Statistical MT Results

(1) Compress the big LM to 676 GB
(2) Decode with 1 TB RAM
(3) Make three WMT 2013 submissions

|  | Czech-English |  | French-English |  | Spanish-English |  |
| ---: | :---: | ---: | :---: | :---: | :---: | ---: |
|  | Rank | BLEU | Rank | BLEU | Rank | BLEU |
| This Work | 1 | 28.16 | 1 | 33.37 | 1 | 32.55 |
| Google | $2-3$ | 27.11 | $2-3$ | 32.62 | 2 | 33.65 |
| Baseline | $3-5$ | 27.38 | $2-3$ | 32.57 | $3-5$ | 31.76 |

## Future Work on Estimation

- Pruning
- Linearly interpolate separately trained models $\rightarrow$ SRI's ARPA output is misleading.
- More smoothing methods
- Parallelization by data splitting


## Outline

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| 1 Estimation fro | 7.1x 0.07 x | ACL 2013 |
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|  |  | NAACL 2013 |

## Raw Queries

Answer language model queries using less time and memory.

$$
\begin{aligned}
& \log p(\text { iran } \mid<\mathrm{s}>\quad)=-3.33437 \\
& \log p(\text { is } \quad \mid<\mathrm{s}>\text { iran } \quad)=-1.05931 \\
& \log p \text { (one } \mid<\mathrm{s}>\text { iran is } \\
& \text { ) }=-1.80743 \\
& \log p(\text { of } \quad \mid<s>\text { iran is one } \quad)=-0.03705 \\
& \log p \text { (the } \quad \text { iran is one of } \quad)=-0.08317 \\
& \log p(\text { few } \quad \text { is one of the })=-1.20788
\end{aligned}
$$

## Example Language Model

Unigrams

| Words | $\log \mathbf{p}$ | $\log \mathbf{b}$ |
| :--- | ---: | ---: |
| $<\mathbf{s}>$ | $-\infty$ | -2.0 |
| iran | -4.1 | -0.8 |
| is | -2.5 | -1.4 |
| one | -3.3 | -0.9 |
| of | -2.5 | -1.1 |

Bigrams
Words $\quad \log p \log b$
<s> iran $-3.3-1.2$
iran is $\quad-1.7 \quad-0.4$
is one $\quad-2.0 \quad-0.9$
one of $\quad-1.4 \quad-0.6$

Trigrams
Words $\quad \log p$
$<\mathrm{s}>$ iran is -1.1
iran is one -2.0
is one of $\quad-0.3$

## Example Queries

## Unigrams

Words $\log p \log b$

| $<\mathrm{s}>$ | $-\infty$ | -2.0 |
| :--- | ---: | ---: |
| iran | -4.1 | -0.8 |
| is | -2.5 | -1.4 |
| one | -3.3 | -0.9 |
| of | -2.5 | -1.1 |

## Query: <s> iran is

$\log p($ is $\mid<\mathrm{s}>$ iran $)=-1.1$

Bigrams
Words $\log p \log b$
<s> iran $-3.3-1.2$
iran is $\quad-1.7 \quad-0.4$
is one $\quad-2.0 \quad-0.9$
one of $\quad-1.4 \quad-0.6$

Trigrams
Words $\quad \log p$
$<\mathrm{s}>$ iran is -1.1
iran is one -2.0
is one of $\quad-0.3$

## Query: iran is of

| $\log p($ of $)$ | -2.5 |
| :--- | ---: |
| $\log b($ is $)$ | -1.4 |
| $\log b($ iran is $)$ | +-0.4 |
| $\log p($ of $\mid$ iran is $)$ | $=-4.3$ |

## Trie Based

CMU-Cambridge Early implementation SRI Popular, considered fast, high-memory
IRST Smaller than SRI, single-threaded
MIT Batch querying
TPT Memory locality
Joshua Java, "not as scalable as the SRILM" [Li et al]
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## Lossy Low-Memory

Rand Bloom maps
Shef Minimal perfect hashing
Google Minimal perfect hashing, larger than Shef

## KenLM Features

- Faster than all baselines
- Lowest lossless memory
- Multithreaded
- Quick loading via memory mapping
- Easier to compile


## Data Structures

Probing Fast.
Trie Small. But still fast.

## Probing

Hash every $n$-gram to a 64-bit integer. Ignore collisions. Store $n$-grams in custom linear probing hash tables.

Fastest, 24 bytes/n-gram (still less than SRI).

## Trie

Reverse $n$-grams, arrange in a trie.


Smaller than most, faster than all but probing.

## Optimizing the Trie

CPU Interpolation search instead of binary search [Yehoshua et al, 1978]

RAM Pack at the bit level i.e. $\log p$ has no sign bit

## Options to Save More Memory

Cluster floats into $2^{f}$ bins, store $f$ bits/float.

Chop bits from integer sequences, store offsets.
[Whittaker and Raj, 2001; Raj and Whittaker, 2003]

## Experiment: Raw Queries

Task Score the English Gigaword corpus<br>Model 5-gram Europarl + deduped news crawl 2011

Queries/ms Excludes loading and file reading time
Loaded RAM Resident after loading
Peak RAM Peak virtual after scoring

Raw Queries: Exact Variants


## Raw Queries: All Tested Variants



## Experiment: Translate 3003 Sentences

Task WMT 2011 French-English baseline Decoder Moses<br>Model 5-gram Europarl+News LM (same as before)<br>Formalism Phrase-based from Europarl

Time Total wall time, including loading
Memory Total resident memory after decoding

## Moses Benchmarks: Single Threaded



## New: Yasuhara et al, EMNLP 2013

"An Efficient Language Model Using Double-Array Structures"

- $19 \%$ less RAM and $4-9 \%$ faster than this work's probing method.
- More RAM than the trie method.
- 4 days to build a data structure with 936 million $n$-grams.

WMT 2013 Adoption: Any Task
This Work 17


## Outline

## Speed RAM Published

## 1 Estimation from Text <br> 2 Raw Queries <br> 3 Decoding <br> 7.1x 0.07x ACL 2013 <br> 2.4x 0.57x WMT 2011 3.2-10.0x $0.85 \times\left\{\begin{array}{l}\text { IWSLT 2011, } \\ \text { EMNLP 2012, } \\ \text { NAACL 2013 }\end{array}\right.$

Decoding performance includes $\approx 1.15 x$ speedup from raw queries.
Baseline: SRILM and cube pruning (more later).

## Parsing-Based MT is Slow

## 26 CPU hours to translate 3000 sentences



French-English system from Ammar et al [2013] using cdec, a 4-gram LM, and cube pruning with beam size 200.

## Decoding Example: Input

Le garçon a vu l'homme avec un télescope

Decoding Example: Parse with SCFG


Decoding Example: Read Target Side


## Decoding Example: One Constituent



| $\overbrace{X: V}^{X: V P}$ | $X: N P$ |
| :---: | :---: |
| a vu | I'homme |
| Hyp | Hyp |
| seen | man |
| saw | the man |
| view | some men |



$X: V P$

a $v u$
Hyp Score

| seen | -3.8 | man | -3.6 |
| :--- | :--- | :--- | :--- |
| saw | -4.0 | the man | -4.3 |
| view | -4.0 | some men | -6.3 |

$X: V P$
a vu l'homme
Hypothesis Score
saw the man -6.9
seen the man
-7.6
saw man -8.3
saw some men -8.5
view man -8.5
seen man -8.8
view the man $\quad-8.9$
seen some men -9.5
view some men -10.8



## Appending Strings

Hypotheses are built by string concatenation. Language model probability changes when this is done:

$$
\frac{p(\text { saw the man })}{p(\text { saw }) p(\text { the } \operatorname{man})}=\frac{p(\text { the } \mid \text { saw }) p(\text { man } \mid \text { saw the })}{p(\text { the })}
$$

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$$

Log probability is part of the score $\Longrightarrow$ Scores do not sum
$\Longrightarrow$ Local decisions may not be globally optimal $\Longrightarrow$ Search is hard.

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1 Estimation from Text<br>2 Raw Queries 7.1x 0.07x ACL 2013<br>2.4x 0.57x WMT 2011<br>3 Decoding 3.2-10.0x $0.85 x$<br>1 State and Recombination<br>2 Score Estimates<br>3 New Search Algorithm<br>EMNLP 2012<br>NAACL 2013

## Appending Strings

Hypotheses are built by string concatenation. Language model probability changes when this is done:
$c($ saw $\bullet$ the man $)=\frac{p(\text { saw the man })}{p(\text { saw }) p(\text { the man })}=\frac{p(\text { the } \mid \operatorname{saw}) p(\text { man } \mid \text { saw the })}{p(\text { the })} p(\operatorname{man} \mid$ the $) \quad$
What words does correction $c$ examine?

## Markov Assumption

A 5-gram language model uses up to 4 words of context: $p($ man $\mid<\mathrm{s}>$ the boy saw the $)=p(\operatorname{man} \mid$ the boy saw the $)$

Correction c examines up to 4 words from each string: $\mathrm{c}\left(<\mathrm{s}>\mathbb{F}^{\text {the boy saw the }- \text { man with a telescope }-1 .) ~}\right.$ Right State

Left State
End of state
End of state

## Hypotheses have State



Right State

| , | 1s | h |
| :---: | :---: | :---: |
| countries that maintain diplomatic -1 | ties | F with North Korea |
| untries that maintain some -1 diplo |  | with North Korea |

## State Controls Recombination



## Left State

 countries that maintain diplomatic -1 relations $\begin{aligned} &- \text { with North Korea . } \\ & \text { ties }\end{aligned}$ countries that maintain some -1 diplomatic ties $\leftarrow$ with North Korea .The decoder may recombine hypotheses with equal state.

## State Controls Recombination



## Left State

 countries that maintain some $\dashv$ diplomatic ties $\forall$ with North Korea .

The decoder may recombine hypotheses with equal state.

Smaller state<br>$\Longrightarrow$ More recombination<br>$\Longrightarrow$ Reason over more hypotheses at once<br>$\Longrightarrow$ Improved time-accuracy tradeoff.

## Efficiently Minimizing State

Li et al [2008] Criterion for state minimization. "Inefficient implementation"

This Work (IWSLT 2011) Repurpose log probability sign bit. Use existing lookups. Encode state to make queries faster.

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Li et al [2008] Criterion for state minimization. "Inefficient implementation"

This Work (IWSLT 2011) Repurpose log probability sign bit. Use existing lookups. Encode state to make queries faster.
$11 \%$ faster than right state minimization alone. On hierarchical Chinese-English with beam size 1000.

## Outline

## Speed RAM

1 Estimation from Text
2 Raw Queries
3 Decoding
1 State and Recombination
2 Score Estimates
3 New Search Algorithm 3.2-10.0x 0.85x

Published 7.1x 0.07x ACL 2013 2.4x 0.57x WMT 2011

IWSLT 2011
EMNLP 2012
NAACL 2013

## Baseline: How to Score a Fragment

| $\log p_{5}($ is $)$ | $=-2.63$ |
| ---: | :--- | ---: |
| $\log p_{5}($ one $\mid$ is $)$ | $=-2.03$ |
| $\log p_{5}($ of $\mid$ is one $)$ | $=-0.24$ |
| $\log p_{5}($ the $\mid$ is one of $)$ | $=-0.47$ |
| $+\log p_{5}($ few $\mid$ is one of the $)$ | $=-1.26$ |
| $=\log p_{5}($ is one of the few $)$ | $=-6.62$ |

# The Problem: Lower Order Entries 

## 5-Gram Model: $\log p_{5}$ (is) $=-2.63$ Unigram Model: $\log p_{1}($ is $)=-2.30$ Same training data.

## The Problem: Lower Order Entries

## 5-Gram Model: $\log p_{5}$ (is) $=-2.63$ <br> Unigram Model: $\log p_{1}$ (is) $=-2.30$

Same training data.


## Build One Model For Each Order

|  | Baseline | Lower <br> $\log p_{5}($ is $)$ |
| ---: | ---: | :--- |
| $=-2.63$ | $-2.30=\log p_{1}$ |  |
| $\log p_{5}($ one $\mid$ is $)$ | $=-2.03$ | $-1.92=\log p_{2}$ |
| $\log p_{5}($ of $\mid$ is one $)$ | $=-0.24$ | $-0.08=\log p_{3}$ |
| $\log p_{5}($ the $\mid$ is one of $)$ | $=-0.47$ | $-0.21=\log p_{4}$ |
| $+\log p_{5}($ few $\mid$ is one of the $)$ | $=-1.26$ | $-1.26=\log p_{5}$ |
| $=\log p_{5}($ is one of the few $)$ | $=-6.62$ | $-5.77=\log p_{\text {Low }}$ |

## Storing Lower Order Models

One extra float per entry, except for longest order. Unigrams
Words $\quad \log p_{5} \log b_{5} \log p_{1}$
australia $\quad-3.9 \quad-0.6 \quad-3.6$
is $\quad-2.6 \quad-1.5 \quad-2.3$
one $\quad-3.4 \quad-1.0 \quad-2.9$
of $\quad-2.5 \quad-1.1 \quad-1.7$

## No need for backoff $b_{1}$

If backoff occurs, use of $p_{5}$ is appropriate.

## Storing Lower Order Models

One extra float per entry, except for longest order. Unigrams
Words $\quad \log p_{5} \log b_{5} \log p_{1}$
australia $\quad-3.9 \quad-0.6 \quad-3.6$
is $\quad-2.6 \quad-1.5 \quad-2.3$
one $\quad-3.4 \quad-1.0 \quad-2.9$
$\begin{array}{llll}\text { of } & -2.5 & -1.1 & -1.7\end{array}$

## No need for backoff $b_{1}$

If backoff occurs, use of $p_{5}$ is appropriate.

Related: store upper bounds [Carter et al, also EMNLP 2012].

So far Better estimates but more memory. Next Worse estimates with less memory.

## Pessimism

## Assume backoff all the way to unigrams.

 $q($ is one of $)=p($ is one of $) b($ is one of $) b($ one of $) b($ of $)$
## Sentence Scores Are Unchanged

$$
q(<\mathrm{s}>\cdots</ \mathrm{s}>)=p(<\mathrm{s}>\cdots</ \mathrm{s}>)
$$

because $b(\cdots</ s>)=1$

## Pessimism

## Assume backoff all the way to unigrams.

 $q($ is one of $)=p($ is one of $) b($ is one of $) b($ one of $) b($ of $)$
## Sentence Scores Are Unchanged

$$
\begin{gathered}
q(<\mathrm{s}>\cdots</ \mathrm{s}>)=p(<\mathrm{s}>\cdots</ \mathrm{s}>) \\
\text { because } b(\cdots</ \mathrm{s}>)=1
\end{gathered}
$$

Telescoping

$$
\begin{aligned}
q(\text { is }) & =p(\text { is }) b(\text { is }) \\
q(\text { one } \mid \text { is }) & =p(\text { one } \mid \text { is }) \frac{b(\text { is one }) b(\text { one })}{b(\text { is })}
\end{aligned}
$$

## Saving Memory

| Unigrams |  |  | Unigrams |  |  |
| :--- | ---: | ---: | :--- | :--- | :--- |
| Words | $\log p_{5}$ | $\log b_{5}$ |  | Words |  |
| $\boldsymbol{\operatorname { l o g }} q_{5}$ |  |  |  |  |  |
| Wustralia | -3.9 | -0.6 |  |  |  |
| is | -2.6 | -1.5 |  | australia | -4.5 |
| one | -3.4 | -1.0 | Compress | is | -4.1 |
| of | -2.5 | -1.1 |  | on | -4.4 |
| of |  | of | -3.6 |  |  |

One less float per entry, except for longest order.

Backoff smoothing with RAM comparable to stupid backoff's counts. Includes Kneser-Ney.

## Experiments

Task WMT 2011 German-English
Decoder Moses with probing LM + state minimization
LM 5-gram from Europarl, news commentary, and news
Grammar Target-syntax and hierarchical systems
Parser Collins

Target-Syntax Model Score


## Hierarchical Model Score



## Hierarchical BLEU



## Memory

Effect of adding or removing a float per entry.

| Structure | Baseline (MB) | Change (MB) | $\mathbf{\%}$ |
| :--- | ---: | ---: | ---: |
| Probing | 4,072 | 517 | $13 \%$ |
| Trie | 2,647 | 506 | $19 \%$ |
| 8-bit quantized trie | 1,236 | 140 | $11 \%$ |
| 8-bit minimal perfect hash | 540 | 140 | $26 \%$ |

## Outline

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| 3 Decoding | $3.2-10.0 \times$ | $0.85 \times$ |

$X: V P$

a $v u$
Hyp Score seen -3.8 saw -4.0 the man -4.3 view -4.0 some men -6.3

Pruning is Approximate

## a vu l'homme

Hypothesis Score
-6.9
-7.6
saw man -8.3
-3.6 sow some men
view man
seen man $\quad-8.8$
view the man -8.9 seed some men 9.5

saw the man
seen the man

Lew some men -10.?

## Beam Search [Lowerre, 1976; Chiang, 2005]

|  | man | $-\mathbf{3 . 6}$ | the man | $-\mathbf{4 . 3}$ | some men | $-\mathbf{6 . 3}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | ---: |
| seen $-\mathbf{3 . 8}$ | seen man -8.8 | seen the man -7.6 | seen some men | -9.5 |  |  |
| saw -4.0 | saw man -8.3 | saw the man | -6.9 | saw some men | -8.5 |  |
| view -4.0 | view man -8.5 | view the man -8.9 | view some men -10.8 |  |  |  |

Baseline: Cube Pruning [Chiang, 2007]

man -3.6 the man -4.3 some men -6.3<br>seen -3.8 Queue<br>saw -4.0<br>view -4.0

## Queue

Hypothesis
Sum
$\rightarrow$ seen man
$-3.8-3.6=-7.4$

## Baseline: Cube Pruning [Chiang, 2007]

|  | man | -3.6 | the man -4.3 | some men |
| :--- | :--- | :--- | :--- | :--- |$-6.3$

## Queue

Hypothesis
$\begin{array}{cl}\rightarrow \text { saw man } & -4.0-3.6=-7.6 \\ \text { seen the man } & -3.8-4.3=-8.1\end{array}$

Sum

## Baseline: Cube Pruning [Chiang, 2007]

|  | man | -3.6 | the man -4.3 | some men | -6.3 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| seen -3.8 | seen man -8.8 | Queue |  |  |  |
| saw -4.0 | saw man -8.3 | Queue |  |  |  |
| view -4.0 | Queue |  |  |  |  |

## Queue

Hypothesis
$\rightarrow$ view man
seen the man
saw the man $\quad-4.0-4.3=-8.3$

Problem With Cube Pruning

Hyp
is a
are a

Hypothesis countries that countries which country

Hypothesis
is a countries that are a countries that are a countries which

No notion that "a countries" is bad.

## Problem With Cube Pruning

Hypothesis countries that is a are a

Hypothesis
is a countries that are a countries that are a countries which

No notion that "a countries" is bad.

## Idea: group by outermost words.

## Example Hypotheses

Left State
Right State countries that -1 maintain diplomatic $\begin{gathered}\text { relations } \\ \text { ties } \\ \text { with North Korea . }\end{gathered}$ countries that have -1 an embassy in $F$ DPR Korea. country -1 that maintains some diplomatic ties $\mathbb{1}$ in North Korea nations which has $\dashv$ some diplomatic ties - with DPR Korea . country -1 that maintains some diplomatic ties with DPR Korea .

## Example Hypotheses

Left State
(countries that (nations which has $\dashv \diamond \vdash$ with DPR Korea .)
(countries that have $\dashv \diamond \vdash$
(country $\quad \dashv \diamond \vdash \quad$ in North Korea .)
(country

Right State $\dashv \diamond \vdash$ with North Korea .) DPR Korea .) $\dashv \diamond \vdash$ with DPR Korea.)
$\dashv$ Left state is completely present.
$\diamond$ Stands for elided words
$\vdash$ Right state is completely present.

## Group by Leftmost Word

## countries


country

## Reveal Common Words in Each Group

(countries that $\diamond$ Korea .)

$(\epsilon \diamond \epsilon) \longrightarrow$ (nations which has $\dashv \diamond \vdash$ with DPR Korea .)

(country $\dashv \diamond$ Korea .)

## Alternate Sides Until Tree is Full

## (countries that $\dashv \diamond \vdash$ with North Korea .)



## Using Rules

$$
\begin{array}{lc}
\text { is a } X: N P 1</ \mathrm{s}> & \\
\text { turns into } & \text { turns into } \\
\text { is a }(\epsilon \diamond \epsilon)</ \mathrm{s}> & \underbrace{(\epsilon \diamond \epsilon)}_{X: V 1} \text { the } \underbrace{(\epsilon \diamond \epsilon)}_{X: N 2}
\end{array}
$$

## Exploring and Backtracking

Does the LM like "is a (countries that $\diamond$ Korea .) $</ \mathrm{s}>$ "?
Yes Try more detail.
No Consider alternatives.

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No Consider alternatives.

Formally: priority queue containing breadcrumbs.

## Split and Leave Breadcrumbs

(countries that $\dashv \diamond \vdash$ with North Korea .)

## (countries that $\diamond$ Korea .)

(countries that have $\dashv \diamond \vdash$ DPR Korea .)
$(\epsilon \diamond \epsilon) \longrightarrow$ (nations which has $\dashv \diamond \vdash$ with DPR Korea .)

(country $\dashv \diamond$ Korea .)


## Split and Leave Breadcrumbs

(countries that $\dashv \diamond \vdash$ with North Korea .)
$\xrightarrow{ }$

## (countries that $\diamond$ Korea .)

(countries that have $\dashv \diamond \vdash$ DPR Korea .)


## The queue entry

$$
\text { is a }(\epsilon \diamond \epsilon)</ \mathrm{s}>
$$

## splits into

# Zeroth Child "is a (countries that $\diamond$ Korea .) </s>" Other Children "is a $(\epsilon \diamond \epsilon)[1+]</ \mathrm{s}>$ " 

Children except the zeroth.

## Summary So Far

A priority queue contains competing entries: is a (countries that $\diamond$ Korea .) $</ \mathrm{s}>$
$(\epsilon \diamond \epsilon)$ the $(\epsilon \diamond \epsilon)$ is a $(\epsilon \diamond \epsilon)[1+]</ \mathrm{s}>$

The algorithm pops the top entry, splits a non-terminal, and pushes.

## Summary So Far

A priority queue contains competing entries: is a (countries that $\diamond$ Korea .) $</ \mathrm{s}>$
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Next: Scoring queue entries

## Scores come from the best descendant:

Score $(\epsilon \diamond \epsilon)=$
Score(countries that $\dashv \diamond \vdash$ with North Korea .)


Score $(\epsilon \diamond \epsilon)[1+]=$
Score(nations which has $\dashv \diamond \vdash$ with DPR Korea .)

## Estimates Update as Words are Revealed

$$
\begin{array}{ll}
\text { is a }(\epsilon \diamond \epsilon)</ \mathrm{s}>\longrightarrow \text { is a (countries that } \diamond \text { Korea } .)</ \mathrm{s}> \\
p(\text { is }) & p(\text { is }) \\
p(\mathrm{a} \mid \text { is }) & p(\mathrm{a} \mid \text { is }) \\
p(\text { countries }) & p(\text { countries } \mid \text { is a) } \\
p \text { (that } \mid \text { countries }) & p \text { (that } \mid \text { is a countries }) \\
p(</ \mathrm{s}>) & p(</ \mathrm{s}>\mid \text { Korea } .)
\end{array}
$$

## Summary: Processing a Constituent

(1) Initialize: Push rules onto a priority queue.
(2) Best-First Loop:

- Pop the top entry.
- If it's complete, add to the beam.

Otherwise, split and push.

- Finalize: Convert the beam to a tree (lazily).


## Summary: Processing a Constituent

(1) Initialize: Push rules onto a priority queue.
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Process constituents in bottom-up order (like cube pruning).

Coarse-to-Fine [Zhang et al, 2008; Petrov et al, 2008]
Decode multiple times, each with more detail:

- LM order
- Word classes

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- LM order
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## Key Difference

Coarse-to-Fine Lock-step refinement This Work Locally refine on demand

## Future Work

- Use this work for each decoding pass
- Word classes for this work


## Exact Algorithms

- Weighted finite state transducers [lglesias et al, 2011]
- Integer linear programming [Rush et al, 2011]
- Later: upper bounds and LM refinement [Aziz et al, WMT 2013]

Currently intractable for large MT (7 hours for a 7-word sentence)
$\Longrightarrow$ Used as first pass of approximate coarse-to-fine.

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## Key Difference

Approximation based on average-case scores before expanding hypotheses.

## Experiment

Task WMT 2011 German-English Built [Koehn et al, 2011]
Model Hierarchical
Decoder Moses
Baseline Queries + State + Rest Costs

## Moses Hierarchical



## Moses Hierarchical





## Summary

Optimized the entire LM pipeline from estimation to search.

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## Comparison

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## Decoder Support


kheafield.com/code

## Questions?

## Full State Minimization

Keep only words that might form cross-hypothesis n-grams.

## Left State [Joshua]

For any word $w$, does the model contain:

$$
\begin{array}{ll}
w \text { countries that maintain some } & x \\
w \text { countries that maintain } & x \\
w \text { countries that }
\end{array}
$$

$\Longrightarrow$ Left state minimizes to "countries that" $\dashv$.

## Full State Minimization

Keep only words that might form cross-hypothesis $n$-grams.

## Left State [Joshua]

For any word $w$, does the model contain:
$w$ countries that maintain some $X$
$w$ countries that maintain $x$
$w$ countries that
$\Longrightarrow$ Left state minimizes to "countries that" $\dashv$.
Right State [SRI, Rand, Joshua]
For any word $w$, does the model contain:
with North Korea . w
$\Longrightarrow$ Right state minimizes to $\vdash$ "with North Korea ."

## Related Work on State

Joshua Left and right but "inefficient implementation" SRI Right only, additional lookups

This Work Repurpose memory, existing lookups Also: encode state to make queries faster

## Experimental Setup

Task NIST Chinese-English [Koehn, 2011]
LM Xinhua and AFP from English Gigaword $4+$ Parallel Data
Grammar Hierarchical
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11\% faster

## Experiments: Systems

Hierarchical with Moses [Koehn, 2012]

- German-English also ported to cdec, Joshua, and Jane
- English-German
- Chinese-English

Target Syntax with Moses [Koehn, 2012]

- German-English
- English-German

Tree-to-Tree with cdec [Ammar et al, 2013]

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## Experiments: Systems and Scenarios

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Baseline and improved rest costs, 2-3 flavors of cube pruning.

